Improved Cryptanalysis of Py

Paul Crowley

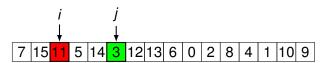
LShift Ltd

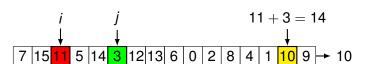
Royal Holloway Information Security Group Seminar, May 2006

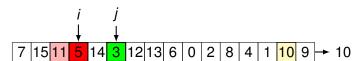
Overview

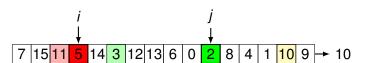
- RC4
- Py
- SPP attack
- Our attack

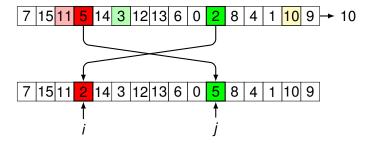
7 | 15 | 11 | 5 | 14 | 3 | 12 | 13 | 6 | 0 | 2 | 8 | 4 | 1 | 10 | 9

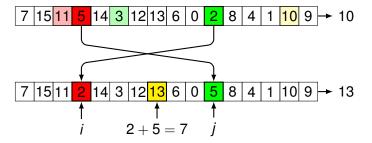


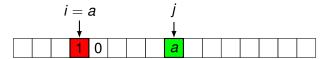


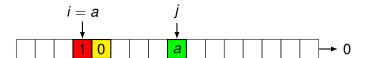


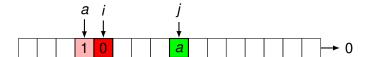


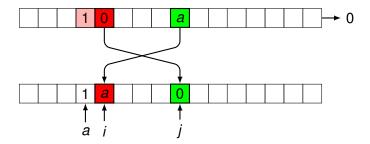


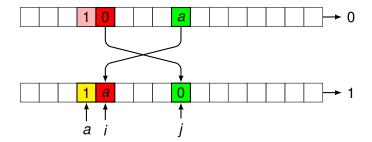








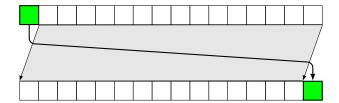




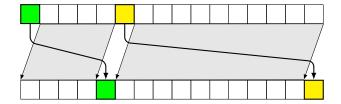
Py

- eSTREAM entrant by Eli Biham and Jennifer Seberry
- Fast in software (2.6 cycles/byte on some platforms)
- SPP attack: 2⁸⁸ bytes of output
- Our attack: 2⁷² bytes

Rolling arrays



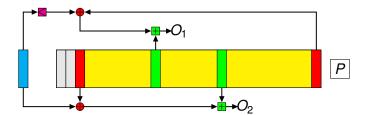
Rolling swaps



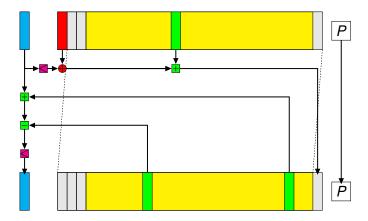
Py internal state



Output



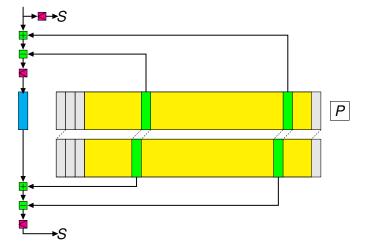
Update



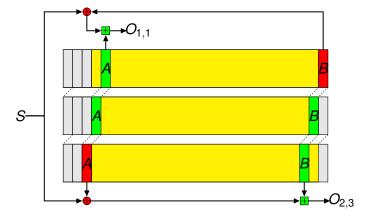
SPP attack

- Gautham Sekar, Souradyuti Paul, Bart Preneel
- Defines event L with $Pr[L] \approx 2^{-41.91}$
- When L occurs, two output bits are the same

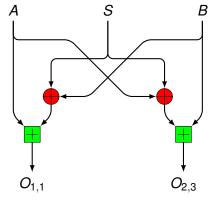
Event L(1)



Event *L* (2)



Result of event L



SPP distinguisher

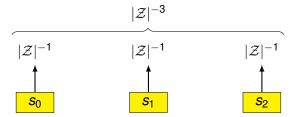
- Examine 2⁸⁵ O_{1,1}, O_{2,3} pairs (ie 2⁸⁸ bytes)
- Count how many pairs have equal low bits
- Report "Py" if above a certain threshold, otherwise "random"
- How do we choose the threshold?

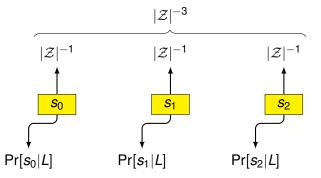
- Thomas Baignères, Pascal Junod, Serge Vaudenay
- Optimal distinguisher chooses the distribution which has the highest probability of producing the observed output
- Neyman-Pearson likelihood ratio test

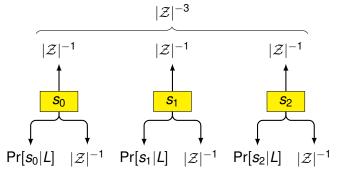
 S_0 S_1 S_2

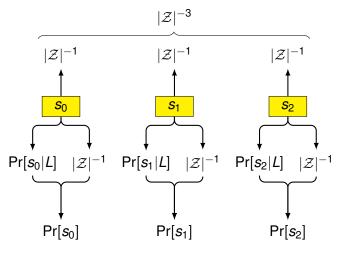


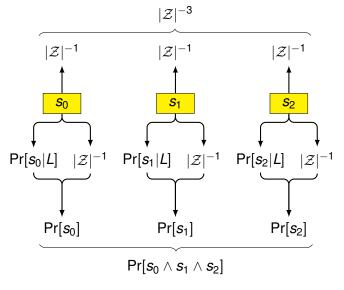












- We score each sample s with $LLR(s) = log(\frac{Pr[s|Py]}{Pr[s|Random]})$
- Sum of scores is log-likelihood ratio for whole sample
- If score is positive, output Py
- Otherwise, output Random

Efficacy of optimal distinguisher

• Where distribution is "close" to uniform random, efficacy

$$\beta = |\mathcal{Z}| \sum_{z \in \mathcal{Z}} \left(\Pr[z] - \frac{1}{|\mathcal{Z}|} \right)^2$$

Efficacy of optimal distinguisher

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- Need around $\frac{2}{\beta}$ samples for advantage $> \frac{1}{2}$

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$$\frac{2}{\beta}$$
 samples for advantage $> \frac{1}{2}$

• If output only biased when event L occurs:

$$\beta = \Pr[L]^2 \left(|\mathcal{Z}| \left(\sum_{z \in \mathcal{Z}} \Pr[z|L]^2 \right) - 1 \right)$$

Efficacy of optimal distinguisher

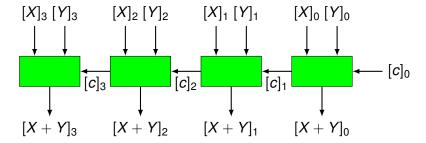
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- Need around $\frac{2}{\beta}$ samples for advantage $> \frac{1}{2}$
- If output only biased when event L occurs: $\beta = \Pr[L]^2 \left(|\mathcal{Z}| \left(\sum_{z \in \mathcal{Z}} \Pr[z|L]^2 \right) 1 \right)$
- SPP attack: $\beta = \Pr[L]^2$ so around 2^{85} samples

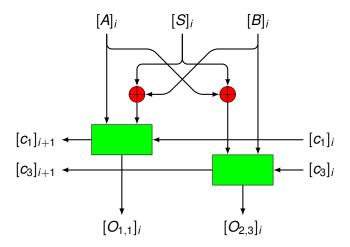
Improving on the attack

- Use all bits of $O_{1,1}, O_{2,3}$
- Group output by column bitwise
- Find exact probability $Pr[O_{1,1}, O_{2,3} = o_{1,1}, o_{2,3}|L]$
- Apply optimal distinguisher

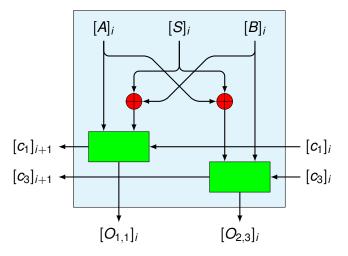
Addition

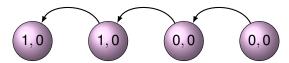


Carry propagation

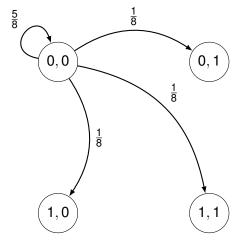


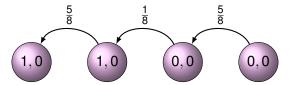
Carry propagation

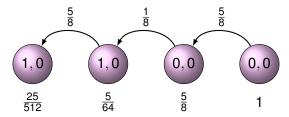


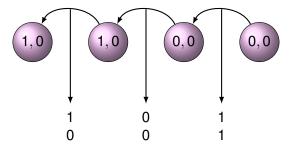


Transition probabilities

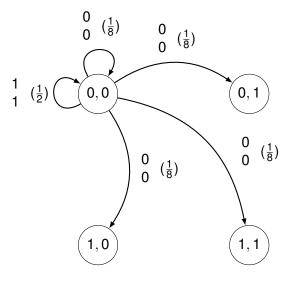


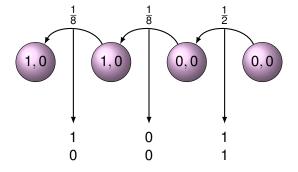


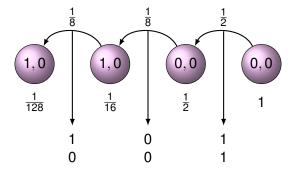


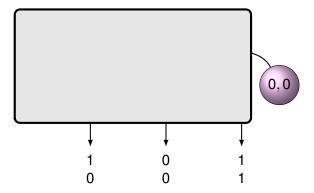


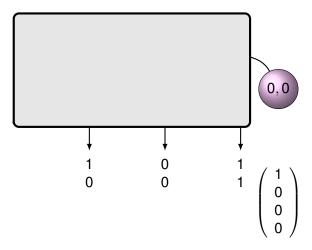
Transition and output probabilities

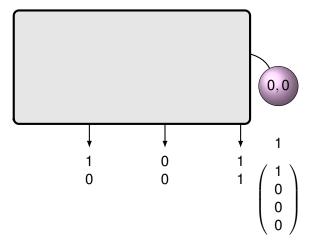


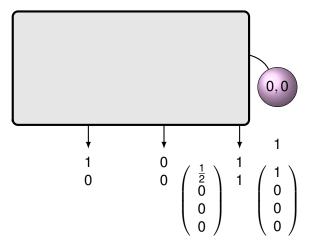


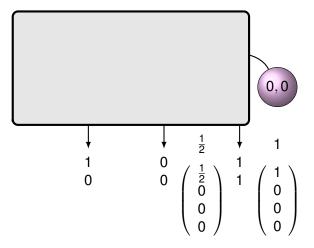


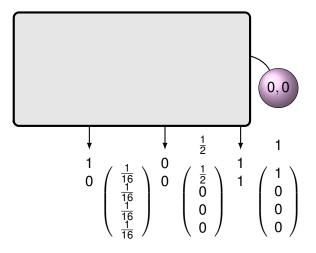


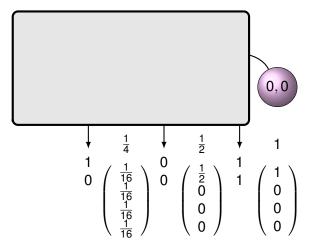


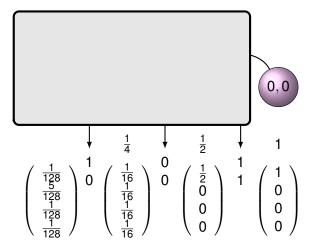


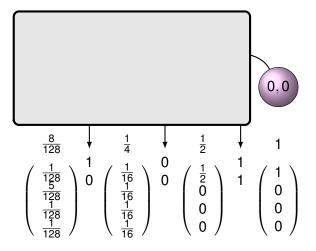


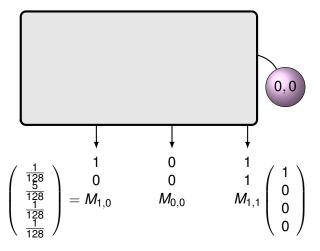












The forward algorithm

$$\Pr\left[\begin{array}{ccc} 1 & 0 & 1 \\ 0 & 0 & 1 \end{array}\right] = \mathbf{1}_{1\times 4} M_{1,0} M_{0,0} M_{1,1} \pi_0$$
 where $\mathbf{1}_{1\times 4} = \left(\begin{array}{ccc} 1 & 1 & 1 & 1 \end{array}\right)$ and $\pi_0 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 0 \end{pmatrix}$

Our attack

- For each sample, use the forward algorithm to find Pr[s|L]
- from which we estimate Pr[s|Py]
- We score each sample s with $LLR(s) = log(\frac{Pr[s|Py]}{Pr[s|Random]})$
- Sum of scores is log-likelihood ratio for whole sample
- If score is positive, output Py
- Otherwise, output Random

$$\sum_{z\in\mathcal{Z}}\Pr[z|L]^2$$

$$\sum_{z \in \mathcal{Z}} \Pr[z|L]^{2}$$

$$= \sum_{z \in \mathcal{Z}} (\mathbf{1}_{1 \times 4} M_{31} M_{30} \dots M_{0} \pi_{0})^{2}$$

$$M_i \in \{M_{0,0}, M_{0,1}, M_{1,0}, M_{1,1}\}$$

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$$= \sum_{z \in \mathcal{Z}} (\mathbf{1}_{1 \times 4} M_{31} M_{30} \dots M_{0} \pi_{0})^{T} M_{0}^{T} \dots M_{30}^{T} M_{31}^{T} \mathbf{1}_{1 \times 4}^{T}$$

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$$= \sum_{z \in \mathcal{Z}} (\mathbf{1}_{1 \times 4} M_{31$$

$$H_i = \sum M_{i-1} M_{i-2} \dots M_1 M_0 \pi_0 \pi_0^T M_0^T M_1^T \dots M_{i-2}^T M_{i-1}^T$$

$$H_{i} = \sum_{T} M_{i-1} M_{i-2} \dots M_{1} M_{0} \pi_{0} \pi_{0}^{T} M_{0}^{T} M_{1}^{T} \dots M_{i-2}^{T} M_{i-1}^{T}$$

$$H_{0} = \pi_{0} \pi_{0}^{T}$$

$$H_{i} = \sum_{M_{i-1}} M_{i-1} M_{i-2} \dots M_{1} M_{0} \pi_{0} \pi_{0}^{T} M_{0}^{T} M_{1}^{T} \dots M_{i-2}^{T} M_{i-1}^{T}$$

$$H_{0} = \pi_{0} \pi_{0}^{T}$$

$$H_{i+1} = \sum_{M \in \{M_{0,0}, M_{0,1}, M_{1,0}, M_{1,1}\}} M H_{i} M^{T}$$

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$$\beta = \Pr[L]^{2} \left(2^{64} \left(\mathbf{1}_{1 \times 4} H_{32} \mathbf{1}_{1 \times 4}^{T} \right) - 1 \right)$$

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$$\approx 60552 \Pr[L]^{2}$$

Conclusions

- We can efficiently calculate the efficacy of HMM-based distinguishers
- Distinguisher advantage is 0.53 given 2⁶⁴ bytes from 2⁸ key/IV pairs
- Advantage is 0.03 given a single 2⁶⁴-byte stream
- Can this be improved still further?

http://www.ciphergoth.org/crypto/py